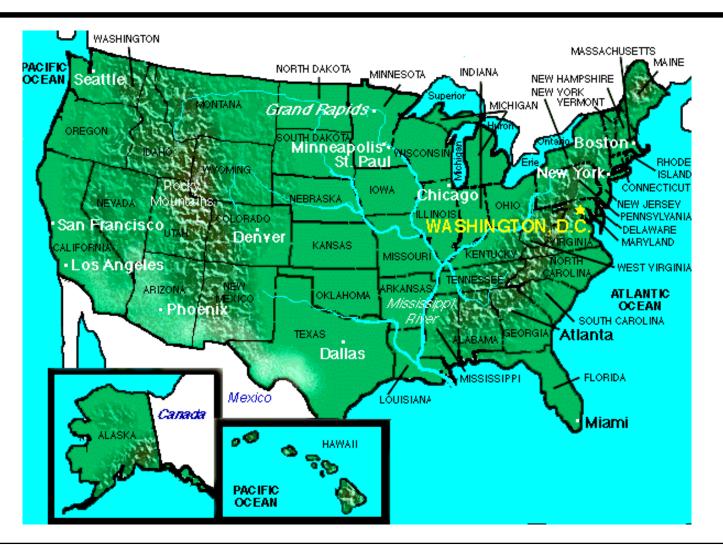
Example: Map coloring



Dynamic backtracking

- $expl_{ik}$
 - explanation for *eliminating* the value v_k for variable x_i
 - set of previously assigned variables that imply that x_i cannot take the value v_k
 - analogous to a combination of mbl_{ik} and conf-set_i generalized to non-binary constraints
 - can be set by forward or backward checking
- assigned
 - sequence of assigned variables
- unassigned
 - set of unassigned variables

Dynamic backtrack: label

```
function dbt-label(i) update\text{-}explanations(i)
if CD_i \neq \{\} then
Select some \ v_k \in CD_i \ \text{ and set } x_i = v_k
push(i, assigned) \ \text{ and } unassigned = unassigned \setminus \{i\}
\mathbf{return} \ (NextVar(), true)
\mathbf{else}
\mathbf{return} \ (i, false)
\mathbf{end} \ dbt-label
```

Updating eliminating explanations

```
procedure update-explanations(i)

for each v_k \in CD_i do

for each completed constraint C that refers to x_i do

if C is not satisfied then

expl_{ik} = vars(C) \setminus \{x_i\}
CD_i = CD_i \setminus \{v_k\}

break inner loop

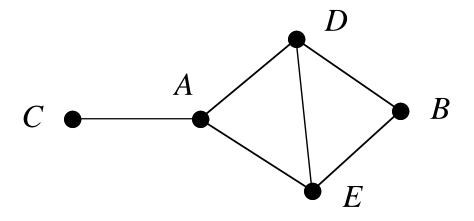
endif

end update-explanations
```

Dynamic backtracking: unlabel

```
function dbt-unlabel(i)
    expl = \bigcup_{r} expl_{ir}
    if expl = \{\} then return (0, false)
    Let j \in expl be the last element pushed onto assigned and
        let v_k be currently assigned to x_i
    for each l assigned after j and each value v_m \in D_l do
        if j \in expl_{lm} then
            expl_{lm} = \{\}
            CD_1 = CD_1 \cup \{v_m\}
        endif
    assigned = assigned \setminus \{j\} and unassigned = unassigned \cup \{j\}
    expl_{ik} = expl \setminus \{j\}
    return (NextVar(), true)
end dbt-unlabel
```

Example



Country	color	red	yellow	blue
A				
B				
C				
D				
E				

Termination of dynamic backtracking

- Setting $expl_{ik} = \{y_1, ..., y_l\}$ is equivalent to asserting the nogood $y_1 = d_1, ..., y_l = d_l \implies x_i \neq v_k$
- Let $\{s_1, ..., s_m\}$ be the variables that $precede\ x_i$ when $expl_{ik}$ is set
- \Rightarrow the following *modified nogood* is entailed

$$s_1 = e_1, ..., s_m = e_m \implies x_i \neq v_k$$

• Lemma: As long as the modified nogood is entailed by some nogood, the algorithm will not visit the node

$$s_1 = e_1, ..., s_m = e_m, x_i = v_k$$

• Corollary: Whenever a nogood is asserted, a previously unentailed modified nogood is entailed

Termination (contd.)

- Lemma: When the algorithm deletes a nogood, the corresponding modified nogood continues to be entailed
- Corollary: The set of entailed modified nogoods grows monotonically
- **Theorem:** Dynamic backtracking terminates